

Assignment 2: Page Table Replacement

- Operation of Page Tables
 - Implement virtual-to-physical address translation and demand paging using a two-level page table
 - Edit methods `allocate_frame` and `*find_physpage` in `pagetable.c`
 - **Paetable**: used to map virtual address and physical addresses. Virtual addresses used by program executed, while physical addresses are used by hardware.
 - **Frame**: chop up space in physical memory in fixed-sized pieces. Each of these frames can contain a single virtual-memory page.
 - **Page**: chop up space in virtual memory in fixed-sized pieces
- Page Replacement: what page to evict to make room
 - Algo
 - **FIFO(First In First Out)**
 - **Clock**(with one ref-bit): evict the oldest page that has not been used
 - One reference bit because one bit is used to keep track if the page is used or not
 - **Exact LRU(Least Recently Used)**: time stamp every reference, evict page with the oldest time stamp
 - **OPT(Optimal Page Replacement)**: also called Belady's algo. Optimal page replacement because lowest fault rate
 - Table columns
 - Hit Rate: displayed after `sim.c` is executed
 - Hit Count: displayed after `sim.c` is executed
 - Miss Count: displayed after `sim.c` is executed
 - Overall Eviction Count: clean eviction count + dirty eviction count
 - Clean Eviction Count: displayed after `sim.c` is executed
 - Dirty Eviction Count: displayed after `sim.c` is executed
 - Code inside respective files and may edit struct frame in `pagetable.h`
- Write up
 - Write in file `README.pdf`
 - Four tables prepared in Task 2 and answer the questions